

Hackable Badge Challenge Walkthrough For SANS EMEA & NCSC UK's CyberThreat24

Solution For "*Echo Service*" By Badge Challenge Author, Secure Impact's Security Engineer, Nathan Taylor

When we start this challenge, we're presented with a screen informing us that USB serial is required, and nothing else. If we connect over serial and then go back and start the challenge we'll see the following:

```
[BOOT] Firmware OK
[BOOT] Complete. Welcome!
Send a single LF to start the challenge.
```

If we send an LF character, as instructed, we see "Service ready" at which point anything we send is returned back to us.

As some of you might already be screaming at your screen, this smells like it's going to be a buffer overflow—and it is. If we type a lot of characters in (around 70), the badge just crashes. Sometimes it'll reboot and other times it'll just freeze.

Unfortunately, this is embedded hacking, not your standard local buffer overflow, so proceeding blind at this point would be particularly challenging. Luckily for us, we have a copy of the firmware already... it's on the badge in your hands!

To communicate with the badge at a lower level, we need to use a piece of software called avrdude. While this tool has many options, we're going to just be using it to communicate with the bootloader.

The bootloader on the hackable badges is Optiboot, running at 115200 baud. This is an Arduinocompatible bootloader, so we can use the following command to dump the firmware:

```
avrdude -V -v -pm1284p -carduino -b115200 -PCOM9 -Uflash:r:flash.bin:r
```

Use avrdude --help to get a break-down of these arguments. My badge was assigned COM9 but yours will likely be different.



This created a flash.bin file containing all 128KiB of the on-chip flash. We don't actually want to analyse all of this. Starting at $1FC00_h$ is the bootloader itself, and there's a large amount of FF_h padding between the main program and the bootloader. By looking at where the FFs start we can truncate the file to $149A8_h$.

It's time to get analysing! I'm going to be using Ghidra, but IDA would probably work just fine. As we have a raw binary file, we're going to need to tell Ghidra what it is. A search for "AVR" should list a few options, but in our case we want the 24-bit variant, compiled with GCC.

Format:	Raw Binary	~ 🕕
Language:	avr8:LE:24:xmega:gcc	
Destination Folder:	CT24:/	
Program Name:	flash_wt.bin	
		Options
	OK Cancel	

Once imported, run analysis then jump to the reset handler at code:2cbe.

Normally when performing binary analysis, you might be used to all of your memory sections automatically loading. This is information that's encoded into the headers of an ELF or PE file that let the operating system handle memory loading. In embedded we get no such luxury. Instead, the reset handler contains some basic routines to setup the processor RAM before the rest of the firmware executes. Exactly what these routines look like depends on the target architecture and the compiler, but they are generally rather standard.

We can usually expect to see at least two loops. One will be loading .data and the other will be zeroing .bss . It's not uncommon to see additional regions being loaded, but in our case here we just have the two.

SECU	RE	×						×	••	×	×
TMDA	CI								×	×	
		I					×	×	X		
									•••		•
									~	~	
			do conv dat								
		code:002cc4 ld e0	do_copy_duc	B17.0xd							
		code:002cc5 a0 e0	ldi	X10.0x0							
		code:002cc6 bl e0	ldi	Xhi, 0x1							
		code:002cc7 ea e6	ldi	Z10,0x6a							
		code:002cc8 fb e3	ldi	Zhi,0x3b							
		code:002cc9 01 e0	ldi	R16,0x1							
		code:002cca 0b bf	out	RAMPZ,R16							
		code:002ccb 02 c0	rjmp	LAB_code_002cce							
			*********	******	*****	****					
			* .data init	loop		*					
			* Copies from	n code:009DB5 to mem:0100 (flas	sh.bin:13B6A)	*					
	1		* Length: 0x0	CCE		*					
			********	*****	*****	* * * *					
			LAB_code_002c	cc	XREF[1]:	code:002cd0(j)					
		code:002ccc 07 90	elpm	R0,Z+=>DAT_code_009db5		= 0Dh					
		code:002ccd 0d 92	st	X+=>DAT_mem_0100,R0		= 0Dh					
			LAR and a sola		VDFF(1)-	anda (002ach (†)					
	1	ander 002anne an 2n	LAB_CODE_0020	Vie Oree	AKEF[1]:	code:002ccb(3)					
		code:002cce ae 3c	cpi	Alo, UXCe							
		code:002cdI bl 07	cpc	ADI, KI/							
	·	code:0020d0 d9 I7	oroc	LAD_CODE_002CCC,ZIIG							

This is the first loop, which is copying our static initialisation data out of flash and into memory. AVR has a word size of 16 bits, so while the copy operation is reading from code:9DB5, that corresponds to an offset of 13B6A_h in our raw flash dump.

	code:002cd1 21 e2	ldi	R18,0x21		
	code:002cd2 ae ec	ldi	Xlo,0xce		
	code:002cd3 bd e0	ldi	Xhi,0xd		
-	code:002cd4 01 c0	rjmp	.do_clear_bss_start		
		.do_clear_bs	s_loop	XREF[1]:	code:002cd8(j)
	code:002cd5 1d 92	st	X+=>DAT_mem_0dce,R1		
•		.do_clear_bs	s_start	XREF[1]:	code:002cd4(j)
	code:002cd6 a6 35	cpi	X10,0x56		
	code:002cd7 b2 07	cpc	Xhi,R18		
i	code:002cd8 el f7	brbc	.do_clear_bss_loop,Zflg		

The loop to zero out .bss is a little simpler as it doesn't need to read initialisation data.

We can now go into the memory map and setup our sections. .text already exists as our imported file however we can now truncate it at the start of the initialisation data. Importantly, also make sure we mark it as read-only as this will help Ghidra with analysis.

Name	Start	Þ.	End	Length	R	W	х	Volatile	Artificial	Overl	Туре		Byte Source
.text	code:0000	000	code:009db4.1	0x13b6a	\checkmark		\checkmark				Default		flash.bin[0x0, 0x13b6a]
.data	mem:0100		mem:0dcd	0xcce	\checkmark	\checkmark					Default	\checkmark	flash.bin[0x13b6a, 0xcce]
.bss	mem:0dce		mem:2155	0x1388	\checkmark	\checkmark					Default	\checkmark	init[0x1388]

After we've setup the memory map, it's worth re-running analysis.

At this point, it's always a good time to check strings! We know we have the "Service ready" string to look for, and sure enough we can find it. There are at this point two important things to notice. The first is the "Challenge solved!" string just below and the second is the lack of any cross-reference to the strings. I don't know how well IDA handles this, but Ghidra is struggling to reconcile addresses that address other memory regions. The load instructions are in code: , but due to how AVR works they implicitly reference mem: .



We have a saving grace though. Accesses to these strings will always be performed using the following two instructions:

```
LDI R22, LOW(ADDRESS)
LDI R23, HIGH(ADDRESS)
```

We can write a small script to assemble these two instructions for any given address and then do a byte search for those four bytes!

```
while True:
    offset = int(input(">"), 16)
    print(
        f"6{(offset >> 0) & 0xF:x} "
        f"e{(offset >> 4) & 0xF:x}\t\t"
        f"ldi\tR22,0x{offset & 0xff:02x}"
    )
    print(
        f"7{(offset >> 8) & 0xF:x} "
        f"e{(offset >> 12) & 0xF:x}\t\t"
        f"ldi\tR23,0x{(offset >> 8) & 0xff:02x}"
    )
```

This is a little into the weeds, but working with AVR always ends up being like this. If we use this tool for address 0C1D it tells us the bytes to search for are going to be 6D E1 7C E0. This has one match:



The function at this address also looks like what we might expect; I've already named a few of these functions for simplicity.





Ghidra's decompiler struggles with AVR quite substantially so it may be easier to follow along in the disassembly instead. The majority of this function is a loop that reads from serial and writes values to the stack, breaking out of the loop when a newline character is received.

P	_loop_head							
	code:008cb9 83 e	a	ldi	R24,0xa3				
	code:008cba 97 e	1	ldi	R25,0x17				
	code:008cbb 0e 9	4 e2 2f	call	HardwareSerial::available				
	code:008cbd 21 e	0	ldi	R18,0x1				
	code:008cbe 89 2	o	or	R24,R25				
	code:008cbf 09 f	4	brbc	LAB_code_008ccl,Zflg				
	code:008cc0 20 e	0	ldi	R18,0x0				
- L -⊳		LAB	_code_008cc1					
	code:008cc1 22 2	3	and	R18,R18				
	code:008cc2 bl f	3	brbs	_loop_head,Zflg				
	code:008cc3 83 e	a	ldi	R24,0xa3				
	code:008cc4 97 e	1	ldi	R25,0x17				
	code:008cc5 0e 9	4 c0 2f	call	HardwareSerial::read				
	code:008cc7 8a 8	3	std	Y+0x2,u8Val				
	code:008cc8 8a 8	1	ldd	u8Val,Y+0x2				
	code:008cc9 8a 3	0	cpi	u8Val,0xa				
	code:008cca 81 f	4	brbc	_not_newline,Zflg				
	code:008ccb 89 8	1	ldd	u8Val,Y+0x1				
	code:008ccc 28 2	£	mov	R18,u8Val				
	code:008ccd 30 e	0	ldi	R19,0x0				
	code:008cce ce 0	1	movw	u8Val,Y				
	code:008ccf 03 9	6	adiw	u8Val,0x3				
	code:008cd0 a9 0	1	movw	R21R20,R19R18				
	code:008cd1 bc 0	1	movw	R23R22,u8Val				
	code:008cd2 83 e	a	ldi	u8Val,0xa3				
	code:008cd3 97 e	1	ldi	u8Val,0x17				
	code:008cd4 0e 9	4 95 31	call	Print::write				
	code:008cd6 83 e	a	ldi	u8Val,0xa3				
	code:008cd7 97 e	1	ldi	u8Val,0x17				
	code:008cd8 0e 9	4 42 68	call	Print::println				
	code:008cda Of c	0	rjmp	_return				
L ▶		_not	t_newline					
	code:008cdb 4a 8	1	ldd	R20,Y+0x2				
	code:008cdc 89 8	1	ldd	u8Val,Y+0x1				
	code:008cdd 91 e	D	ldi	u8Val,0x1				
	code:008cde 98 0	£	add	u8Val,u8Val				
	code:008cdf 99 8	3	std	Y+0x1,u8Val				
	code:008ce0 88 2	£	mov	u8Val,u8Val				
	code:008cel 90 e	0	ldi	u8Val,0x0				
	code:008ce2 9e 0	1	movw	R19R18,Y				
	code:008ce3 2d 5	£	subi	R18,0xfd				
	code:008ce4 3f 4	£	sbci	R19,0xff				
	code:008ce5 82 0	£	add	u8Val,R18				
	code:008ce6 93 1	£	adc	u8Val,R19				
	code:008ce7 fc 0	1	movw	Z,u8Val				
	code:008ce8 40 8	3	st	Z,R20				
	code:008ce9 cf c	£	rjmp	_loop_head				
└>		_ret	turn					
	code:008cea_ce_5	•	subi	Ylo, 0xbe				

Checking the start of this function, we can see where the stack is initialised. 66 bytes are being allocated on the stack, which in this instance corresponds to a 64 byte buffer and 2 bytes for the buffer index.



	undefined ech	oServiceInner()
undefined	R24:1	<return></return>
undefined2	R25R24:2	u8Val
	echoServiceIn	ner
code:008cad cf 93	push	Ylo
code:008cae df 93	push	Yhi
code:008caf cd b7	in	Ylo,SPL
code:008cb0 de b7	in	Yhi, SPH
code:008cb1 c2 54	subi	Ylo,0x42
code:008cb2 d1 09	sbc	Yhi,Rl
code:008cb3 Of b6	in	R0, SREG
code:008cb4 f8 94	cli	
code:008cb5 de bf	out	SPH,Yhi
code:008cb6 Of be	out	SREG,R0
code:008cb7 cd bf	out	SPL,Ylo
code:008cb8 19 82	std	Y+0x1,R1

Now's the time to pause reading and try and completely reverse engineer this function by hand, if you want. For the rest of us, here's the original source code:

```
static void echoServiceInner() {
    uint8_t iBuffer = 0;
    char aBuffer[64];
    while (1) {
        if (Serial.available()) {
            uint8_t u8Val = Serial.read();
            if (u8Val == '\n') {
                Serial.write(aBuffer, iBuffer);
                Serial.println();
                break;
            }
            aBuffer[iBuffer++] = u8Val;
        }
    }
    // ~~oooooo~~~~ I wonder where this will take us!
    return;
}
```

As the comment there might suggest, our objective is going to be to overwrite the return pointer on the stack. We know we have 66 bytes of allocated stack to clobber, so our payload is going to start with 66 nonsense characters. The next two bytes on the stack are the return address, and then finally we're going to need to include a newline character to trigger the break condition.

The question would be, where do we need to return? Remember that "Challenge solved" string from earlier? Let's go follow that. Using our same script from earlier, address 0C32 will be loaded by the sequence 62 E3 7C E0.



As with last time, there's only a single hit for this sequence. This makes our target return address code: 8D59.



Putting all of that together, we get a payload of

Sending this to the badge, we can see

The challenge was solved, and then the badge crashed and rebooted!

Keep that firmware image loaded in Ghidra; we're going to need it again for the next challenge too.